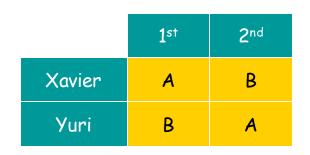
Understanding the Solution

Q. For a given problem instance, there may be several stable matchings. Do all executions of Gale-Shapley yield the same stable matching? If so, which one?

An instance with two stable matchings:

- A-X, B-Y.
- A-Y, B-X.



	1 st	2 nd
Amy	У	X
Brenda	Х	У

Man Optimal Assignments

Definition: Man m is a valid partner of woman w if there exists some stable matching in which they are matched.

Man-optimal assignment: Each man receives the best valid partner (according to his preferences).

- Is man-optimal assignment a matching?
- Simultaneously best for each and every man.

Claim: All executions of GS yield a man-optimal matching, which is a stable matching!

Man Optimality

Claim: GS matching S* is man-optimal.

Proof: (by contradiction)

Suppose some man is paired with someone other than his best partner. Men propose in decreasing order of preference \Rightarrow some man is rejected by a valid partner.

Let **m** be the man who is the first such rejection, and let **w** be the women who is first valid partner that rejects him.

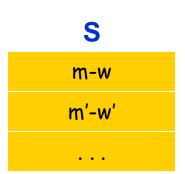
Let **S** be a stable matching where **w** and **m** are matched. In building **S***, when **m** is rejected, **w** forms (or reaffirms) engagement with a man, say **m**', whom she prefers to **m**.

Let w' be the partner of m' in S.

In building **S***, **m**' is not rejected by any valid partner at the point when **m** is rejected by **w**. Thus, **m**' prefers **w** to **w**'.

But w prefers m' to m. Thus w-m' is unstable in S.

since this is the **first** rejection by a valid partner



Man Optimality Summary

Man-optimality: In version of GS where men propose, each man receives the best valid partner.

w is a valid partner of **m** if there exist some stable matching where **m** and **w** are paired