

# Understanding the Solution

Q. For a given problem instance, there may be several stable matchings. Do all executions of Gale-Shapley yield the same stable matching? If so, which one?

An instance with two stable matchings:

- A-X, B-Y.
- A-Y, B-X.

	1 <sup>st</sup>	2 <sup>nd</sup>
Xavier	A	B
Yuri	B	A

	1 <sup>st</sup>	2 <sup>nd</sup>
Amy	Y	X
Brenda	X	Y

# Man Optimal Assignments

**Definition:** Man  $m$  is a **valid partner** of woman  $w$  if there exists some stable matching in which they are matched.

**Man-optimal assignment:** Each man receives the **best** valid partner (according to his preferences).

- Is man-optimal assignment a matching?
- Simultaneously best for each and every man.

**Claim:** **All** executions of GS yield a man-optimal matching, which is a stable matching!

# Man Optimality

S

m-w

m'-w'

...

**Claim:** GS matching  $S^*$  is man-optimal.

**Proof:** (by contradiction)

Suppose some man is paired with someone other than his best partner. Men propose in decreasing order of preference  $\Rightarrow$  some man is rejected by a valid partner.

Let  $m$  be the man who is the **first** such rejection, and let  $w$  be the woman who is **first** valid partner that rejects him.

Let  $S$  be a stable matching where  $w$  and  $m$  are matched.

In building  $S^*$ , when  $m$  is rejected,  $w$  forms (or reaffirms) engagement with a man, say  $m'$ , whom she prefers to  $m$ .

Let  $w'$  be the partner of  $m'$  in  $S$ .

In building  $S^*$ ,  $m'$  is not rejected by any valid partner at the point when  $m$  is rejected by  $w$ . Thus,  $m'$  prefers  $w$  to  $w'$ .

But  $w$  prefers  $m'$  to  $m$ .

Thus  $w$ - $m'$  is unstable in  $S$ .

since this is the **first** rejection  
by a valid partner

# Man Optimality Summary

**Man-optimality:** In version of GS where men propose, each man receives the best **valid** partner.

**w** is a valid partner of **m** if there exist some stable matching where **m** and **w** are paired