CS 401: Computer Algorithm I

DFS / Topological Sort

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Homework 1

Homework 1 will be out later today (Due Feb 18)

- Writing homework
- The first 5 questions are for all the students
- Question 6 is for graduate student only (Undergraduate students who work on Question 6 receive at most 3 bonus points)
- Submit your homework to gradescope

Guidelines:

- You can collaborate, but you must write solutions on your own
- Your solution should be clear, well-organized, and concise. Spell out main idea.
- Sanity Check: Make sure you use assumptions of the problem
- You can use AI/online tools, but you cannot copy AI/online solutions.
- Late homework will be penalized at a rate of 10% of the initial grade per late day (e.g. if your homework receives 80 points and you are late for 4 days, you get 48 points)
- Correctness proofs of ALGORITHMS are NOT REQUIRED

Graph Traversal

Walk (via edges) from a fixed starting vertex *s* to all vertices reachable from *s*.

- Breadth First Search (BFS): Order nodes in successive layers based on distance from *s*
- Depth First Search (DFS): More natural approach for exploring a maze;

Applications of BFS:

- Finding shortest path for unit-length graphs
- Finding connected components of a graph
- Testing bipartiteness

Depth First Search

Follow the first path you find as far as you can go; back up to last unexplored edge when you reach a dead end, then go as far you can



Naturally implemented using recursive calls or a stack

DFS(s) – Recursive version

Initialization: mark all vertices undiscovered

DFS(v) Mark v discovered

> for each edge {*v*, *x*} if (*x* is undiscovered) DFS(*x*)

Mark *v* fully-explored

































































Properties of (undirected) DFS

Like BFS(s):

- DFS(s) visits x iff there is a path in G from s to x
 So, we can use DFS to find connected components
- Edges into then-undiscovered vertices define a tree the "DFS tree" of G

Unlike the BFS tree:

- The DFS tree isn't minimum depth
- Its levels don't reflect min distance from the root
- Non-tree edges never join vertices on the same or adjacent levels

Non-Tree Edges in DFS

BFS tree \neq DFS tree, but, as with BFS, DFS has found a tree in the graph s.t. non-tree edges are "simple" in some way.

All non-tree edges join a vertex and one of its descendents/ancestors in the DFS tree



Directed Graph and Topological Ordering

Directed Graphs



Directed Graphs

Node = intersection, edge = one-way street



Precedence Constraints

In a directed graph, an edge (i, j) means task *i* must occur before task *j*.

Applications

Course prerequisite:

course *i* must be taken before *j*

• Compilation:

must compile module *i* before *j*

• Computing overflow:

output of job i is part of input to job j

 Manufacturing or assembly: sand it before paint it



Directed Acyclic Graphs (DAG)

Def: A DAG is a directed acyclic graph, i.e., one that contains no directed cycles.

Def: A topological order of a directed graph G = (V, E) is an ordering of its nodes as $v_1, v_2, ..., v_n$ so that for every edge (v_i, v_j) we have i < j.



DAGs: A Sufficient Condition

Lemma: If *G* has a topological order, then *G* is a DAG.

Proof. (by contradiction)

Suppose that G has a topological order 1, 2, ..., n and that G also has a directed cycle C.

Let *i* be the lowest-indexed node in *C*, and let *j* be the node just before *i*; thus (j,i) is an (directed) edge.

By our choice of *i*, we have i < j.

On the other hand, since (j,i) is an edge and 1, ..., n is a topological order, we must have j < i, a contradiction

the directed cycle C

DAGs: A Sufficient Condition



Every DAG has a source node

Lemma: If *G* is a DAG, then *G* has a node with no incoming edges (i.e., a source).

Proof. (by contradiction)

Suppose that *G* is a DAG and it has no source

Pick any node v, and begin following edges backward from v. Since v has at least one incoming edge (u, v) we can walk backward to u.

Then, since u has at least one incoming edge (x, u), we can walk backward to x.

Repeat until we visit a node, say w, twice.

Let C be the sequence of nodes encountered between successive visits to w. C is a cycle.



DAG => Topological Order

Lemma: If G is a DAG, then G has a topological order

Proof. (by induction on n)

Base case: true if n = 1.

Hypothesis: Every DAG with n - 1 vertices has a topological ordering.

Inductive Step: Given DAG with n > 1 nodes, find a source node v.

 $G - \{v\}$ is a DAG, since deleting v cannot create cycles.

Reminder: Always remove vertices/edges to use hypothesis

By hypothesis, $G - \{v\}$ has a topological ordering.

Place v first in topological ordering; then append nodes of $G - \{v\}$ in topological order. This is valid since v has no incoming edges.

A Characterization of DAGs

G has a topological order



G is a DAG

Topological Order Algorithm 1: Example



Topological Order Algorithm 1: Example



Topological order: 1, 2, 3, 4, 5, 6, 7

Running time: O(n+m)

- Adjacency list
- Maintain # outgoing edge for each node