

CS151 Fall 2014
Lecture 6 – 9/11
Propositional Logic

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Translating Mathematical Theorem

WEAK: EVERY ODD NUMBER GREATER THAN 5 IS THE SUM OF THREE PRIMES

STRONG: EVERY EVEN NUMBER GREATER THAN 2 IS THE SUM OF TWO PRIMES

VERY WEAK: EVERY NUMBER GREATER THAN 7 IS THE SUM OF TWO OTHER NUMBERS

VERY STRONG: EVERY ODD NUMBER IS PRIME

EXTREMELY WEAK: NUMBERS JUST KEEP GOING

EXTREMELY STRONG: THERE ARE NO NUMBERS ABOVE 7

GOLDBACH CONJECTURES

<http://xkcd.com/1310/>

$$\forall \text{Even}(x), x > 2 \exists \text{Prime}(y), \text{Prime}(z) \quad x = y + z$$

Programs and predicates

Free and bound variables

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function (n) {
  for i = 1 to n do
    print i
}
```

$P(n) :=$ "function(n) prints integer numbers"

$\exists n \quad P(n)$

$Q(n) :=$ "function(n) prints numbers less than 10"

Negations of Quantified Statements

Everyone likes football.

What is the negation of this statement?

Not everyone likes football = There exists someone who doesn't like football.

$$\neg \forall x P(x) \equiv \exists x \neg P(x)$$

(generalized) DeMorgan's Law Say the domain has only three values.

$$\neg \forall x P(x) \equiv \neg(P(1) \wedge P(2) \wedge P(3))$$

$$\equiv \neg P(1) \vee \neg P(2) \vee \neg P(3)$$

$$\equiv \exists x \neg P(x)$$

Negations of Quantified Statements

There is a plant that can fly.

What is the negation of this statement?

Not exists a plant that can fly = every plant cannot fly.

$$\neg \exists x P(x) \equiv \forall x \neg P(x)$$

(generalized) DeMorgan's Law

Say the domain has only three values.

$$\begin{aligned} \neg \exists x P(x) &\equiv \neg(P(1) \vee P(2) \vee P(3)) \\ &\equiv \neg P(1) \wedge \neg P(2) \wedge \neg P(3) \\ &\equiv \forall x \neg P(x) \end{aligned}$$

Order of Quantifiers

There is an anti-virus program killing every computer virus.

How to interpret this sentence?

For every computer virus, there is an anti-virus program that kills it.

$$\forall V \exists P, \text{kill}(P, V)$$

<http://home.mcafee.com/virusinfo/VirusRemovalTools.aspx>

Virus Name	Removal Tool
▶ Sasser	▶ McAfee Stinger
▶ Bagle	▶ McAfee Stinger
▶ Zafi	▶ McAfee Stinger
▶ Mydoom	▶ McAfee Stinger
▶ Lovsan/Blaster	▶ McAfee Stinger
▶ Klez	▶ Klez Removal Tool
▶ Bugbear	▶ Bugbear Removal Tool

"Is Your PC Infected? Don't Worry, We'll Fix It!"

Order of Quantifiers

There is an anti-virus program killing every computer virus.

How to interpret this sentence?

There is one single anti-virus program that kills all computer viruses.

$$\exists P \forall V, \text{kill}(P, V)$$

I have *one* defense good against every attack.

Example: P is ??,
protects against ALL viruses

That's much better!

Order of quantifiers is very important!

More Negations

There is an anti-virus program killing every computer virus.

$$\exists P \forall V, \text{kill}(P, V)$$

What is the negation of this sentence?

$$\begin{aligned} &\neg(\exists P \forall V, \text{kill}(P, V)) \\ &\equiv \forall P \neg(\forall V, \text{kill}(P, V)) \\ &\equiv \forall P \exists V \neg \text{kill}(P, V) \end{aligned}$$

For every program, there is some virus that it can not kill.

Predicate Calculus Validity

Propositional validity

$$p \vee \neg p$$

True *no matter what* the truth value of p is

Predicate calculus validity

$$\forall x \in S [P(x) \vee \neg P(x)]$$

True *no matter what*

- the Domain is
(integers, people, games)
- or the predicates are
($x > 42$, x has red hair, x is similar to Minecraft)

That is, logically correct, independent of the specific content.

Predicate Calculus Validity

Propositional validity

$$(A \rightarrow B) \vee (B \rightarrow A)$$

True *no matter what* the truth values of A and B are

Predicate calculus validity

$$\forall x [[P(x) \rightarrow Q(x)] \vee [Q(x) \rightarrow P(x)]]$$

A fully quantified expression φ of predicate logic is a **theorem** if and only if φ is true for every possible meaning of each of the predicates of φ .

Theorem or not?

$$[\forall x \in \mathbb{Z} P(x)] \vee [\forall x \in (\mathbb{Z}) \neg P(x)]$$

(Not: consider $P(x)$ = "x is a prime")

There's no algorithm that's guaranteed to figure out whether a given fully quantified expression φ is a theorem!
(Goedel's incompleteness theorem, Turing's undecidability of the halting problem)

Arguments with Quantified Statements

Universal instantiation: $\forall x, P(x)$
 $\therefore P(a) \quad [\forall x P(x)] \rightarrow P(a)$

Universal modus ponens: $\forall x, P(x) \rightarrow Q(x)$
 $P(a)$
 $\therefore Q(a)$
 $[[\forall x P(x) \rightarrow q(x)] \wedge P(a)] \rightarrow Q(a)$

Universal modus tollens: $\forall x, P(x) \rightarrow Q(x)$
 $\neg Q(a)$
 $\therefore \neg P(a)$

Universal Generalization

valid rule
$$\frac{A \rightarrow R(c)}{A \rightarrow \forall x.R(x)}$$

providing c is independent of A

e.g. given any number c , $2c$ is an even number

\Rightarrow for all x , $2x$ is an even number.