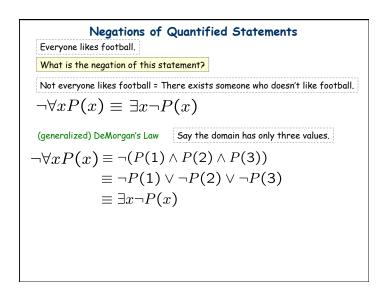


Programs and predicates Free and bound variables function (n) { for $i = 1 \text{ to } n \text{ do } \\ print i } P(n) := "function(n) prints integer numbers" <math>P(n) = P(n)$



Negations of Quantified Statements

There is a plant that can fly.

What is the negation of this statement?

Not exists a plant that can fly = every plant cannot fly.

$$\neg \exists x P(x) \equiv \forall x \neg P(x)$$

(generalized) DeMorgan's Law Say the domain has only three values.

$$\neg \exists x P(x) \equiv \neg (P(1) \lor P(2) \lor P(3))$$
$$\equiv \neg P(1) \land \neg P(2) \land \neg P(3)$$
$$\equiv \forall x \neg P(x)$$

Order of Quantifiers

There is an anti-virus program killing every computer virus.

How to interpret this sentence?

There is one single anti-virus program that kills all computer viruses.

$$\exists P \ \forall V, kill(P, V)$$

I have *one* defense good against every attack.

Example: P is ??. protects against ALL viruses

That's much better!

Order of quantifiers is very important!

Order of Quantifiers

There is an anti-virus program killing every computer virus.

How to interpret this sentence?

For every computer virus, there is an anti-virus program that kills it.

$$\forall V \; \exists P, \mathsf{kill}(P, V)$$

http://home.mcafee.com/virusinfo/VirusRemovalTools.aspx

Virus Name	Removal Tool	
▶ Sasser	► McAfee Stinger	
▶ Bagle	► McAfee Stinger	
▶ Zafi	► McAfee Stinger	
► Mydoom	► McAfee Stinger	
▶ Lovsan/Blaster	► McAfee Stinger	
▶ Klez	▶ Klez Removal Tool	
▶ Bugbear	▶ Bugbear Removal Tool	

"Is Your PC Infected? Don't Worry, We'll Fix It!"

More Negations

There is an anti-virus program killing every computer virus.

 $\exists P \ \forall V, \mathsf{kill}(P,V)$ What is the negation of this sentence?

 $\neg(\exists P \ \forall V, \mathsf{kill}(P, V))$

 $\equiv \forall P \neg (\forall V, \mathsf{kill}(P, V))$

 $\equiv \forall P \exists V \neg \mathsf{kill}(P, V))$

For every program, there is some virus that it can not kill.

Predicate Calculus Validity

Propositional validity

$$p \vee \neg p$$

True *no matter what* the truth value of **p** is Predicate calculus validity

$$\forall x \in S \quad [P(x) \lor \neg P(x)]$$

True no matter what

- the Domain is (integers, people, games)
- or the predicates are (x > 42, x has red hair, x is similar to Minecraft)

That is, logically correct, independent of the specific content.

Predicate Calculus Validity

Propositional validity

$$(A \rightarrow B) \lor (B \rightarrow A)$$

True *no matter what* the truth values of *A* and *B* are Predicate calculus validity

$$\forall x \quad [[P(x) \to Q(x)] \lor [Q(x) \to P(x)]]$$

A fully quantified expression ϕ of predicate logic is a theorem $\it if$ and only if ϕ is true for every possible meaning of each of the predicates of $\phi.$

Theorem or not?

$$[\forall x \in \mathbb{Z} \quad P(x)] \vee [\forall x \in (Z) \neg P(x)]$$

(Not: consider P(x) = x is a prime")

There's no algorithm that's guaranteed to figure out whether a given fully quantified expression $\boldsymbol{\phi}$ is a theorem!

(Goedel's incompleteness theorem, Turing's undecidability of the halting problem)

Arguments with Quantified Statements

Universal instantiation: $\begin{array}{c} \forall x, P(x) \\ \vdots P(a) \end{array} \qquad [\forall x \quad P(x)] \rightarrow P(a)$

 $\forall x, P(x) \rightarrow Q(x)$

Universal modus ponens: P(a) . . Q(a)

 $[[\forall x \quad P(x) \to q(x)] \land P(a)] \to Q(a)$

 $\forall x, P(x) \to Q(x)$

Universal modus tollens: $\neg Q(a)$

 $\therefore \neg P(a)$

Universal Generalization

valid rule
$$\frac{A \to R(c)}{A \to \forall x. R(x)}$$

providing c is independent of A

e.g. given any number c, 2c is an even number

=> for all x, 2x is an even number.