RDF-based Metadata Management in Peer-to-Peer Systems

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Abstract. Peer-to-Peer (P2P) data integration combines the P2P infrastructure with traditional schema-based data integration techniques. One of the primary problems in this research area is the automation of the process of interoperation among heterogeneous data sources. To facilitate this process, metadata are used to add rich semantics to the data sources. In this paper, we propose an RDF-based layered approach to data interoperation in a P2P environment. We use RDF and its vocabulary language RDFS for representation of data and metadata, which are distributed in different peers. For the representation of the mappings of metadata from peer to peer (called *P2P mappings*), we propose a mapping language, which is an RDF-based meta-ontology (called *RDFMS*) that enables metadata management in a decentralized way. We also discuss a thesaurus-based mapping process and query translation in the proposed framework, which aim to provide an efficient way to P2P metadata and data management, respectively.

1 Introduction

Recently, Peer-to-Peer (P2P) techniques are becoming more and more popular in various applications, for its properties such as decentralization, autonomy, and scalability [15]. Typical examples include Napster (www.napster.com) and Gnutella (www.gnutella.com) for file sharing, SETI@home (setiathome.ssl.berkeley.edu) and Grid (www.grid.org) for distributed computing, and AOL instance messaging (www.aol.com) for collaborative work. As an emerging application area, P2P data integration combines the P2P infrastructure with traditional schema-based data integration techniques [1, 3, 6, 7, 11, 16, 17]. For this purpose, one has to first solve the problem of interoperation among heterogeneous data sources. To facilitate automation of the process of data interoperation, metadata (typically in a uniform representation) are used to add rich semantics into data sources. Furthermore, a powerful mechanism is needed for expressing and managing the metadata to support complex queries over the resources in the network.

The example in Figure 1 illustrates the problem more intuitively. The three schema-based data sources (residing in different peers) are heterogeneous in both

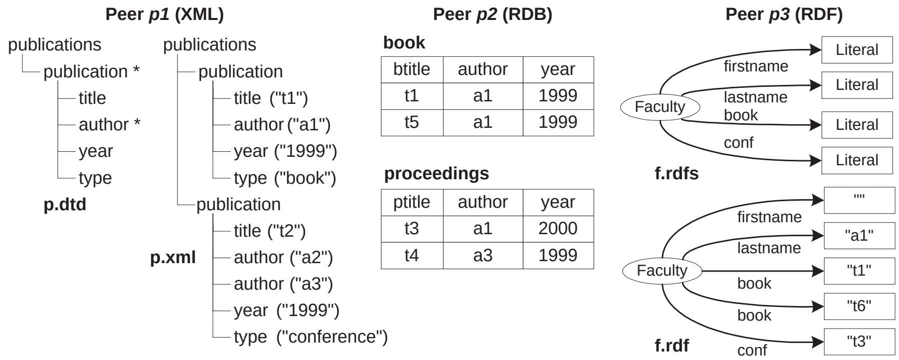

Fig. 1. A motivating example of P2P data interoperation.

syntax and structure. The data in Peer *p1* contains an XML schema in DTD (`p.dtd`) and its data (`p.xml`), where the strings in parentheses are element or attribute values. Peer *p2* is a relational database (RDB) with two relations (`book` and `proceedings`). Peer *p3* contains RDF data (`f.rdf`) with its RDF schema (`f.rdfs`) defined in RDF Schema[1] (RDFS, a vocabulary language for RDF). One notable structural difference among these data sources is that the semantically equivalent terms are formulated in different forms. That is, the two types of publications `book` and `proceedings` are designed as values (instances) in *p1*, as relation names in *p2*, and as property names in *p3*.

The metadata for each data source in the example is a schema, which defines the structure of the data. To enable data access across these data sources, semantic mappings have to be established to hide the syntactic and structural differences. As an example, our previous work has discussed this problem and provided for it a *hybrid* P2P architecture, in which a *super peer* is used as a central control point [10]. In this paper, we propose an RDF-based layered approach to data interoperation in a *pure* P2P environment, in which no central points exist and all peers are in the same roles. More specifically, we use RDF and RDFS for representation of metadata and data, which are distributed in different peers. For the representation of the mappings among metadata in the peers (called *P2P mappings*), we propose a mapping language that is an RDF-based meta-ontology, *RDFMS* (RDF Mapping Schema), which enables metadata management in a decentralized way. We also discuss a thesaurus-based mapping process and query translation in our proposed framework, which aim to provide an efficient and effective way to P2P metadata and data management, respectively.

We base our approach on RDF (through the use of RDFS and RDFMS) because of its particular characteristics as a mechanism for Web resource description. First of all, RDFS organizes metadata declaratively and conceptually (i.e., independently of the document structure) in terms of classes, properties,

[1] http://www.w3.org/TR/rdf-schema/

and data types. Using the metadata, the semantics of concepts and semantic relationships between concepts can be explicitly expressed, thus providing a good foundation for reasoning on the subsequent metadata mapping and query reformulation. Second, RDF properties are actually binary relations between concepts. This feature supports well the representation of P2P mappings, which in our approach are defined as correspondences (other than queries) between metadata in different peers. Finally, a feature of RDF is its extensibility, which allows for easily adding new constructs (e.g., a new mapping type) into the mapping language RDFMS.

The rest of the paper is organized as follows. In Section 2 we describe exiting related work. The layered architecture is introduced in Section 3. In Section 4 we discuss metadata representation in RDF, to which XML and relational schemas are converted. We also discuss in detail the P2P mapping process and mapping representation using RDFMS. The process of query translation based on the P2P mappings is given in Section 5. Finally, Section 6 concludes and discusses future work.

2 Related Work

There are already a number of P2P data management systems or frameworks addressing metadata management in the P2P context.

AXML (ActiveXML) [1] proposes a framework that harnesses *web services* for P2P data integration. In AXML, the mapping process and query processing among peers are respectively realized by *embedding* calls to web services in an XML document and by *firing* these web service calls. In comparison with the schema-based approach proposed in this paper, AXML is approaching the goal of P2P integration at the data level (i.e., XML document centric).

LRM (Local Relational Model) [7], **Hyperion** [3], and **Piazza** [11] provide solutions to P2P metadata and data management, by using *domain relations*, *mapping tables* or *expressions*, and an XQuery-based mapping language, respectively. Compared to these approaches, we believe the RDFMS mapping language proposed in this paper has the advantages of maintainability (via its meta-ontology for mappings representation), extensibility (via the reification property of RDF), inference (via declarative and explicit semantics of RDF), and support for query processing (through RQL, which can even query mappings).

Edutella [16] and **PeerDB** [17] implement metadata and data management in a dynamic (i.e., run-time) way. That is, there are no established mappings between peers at design time. Instead, Edutella uses a Datalog-based query exchange language serving as a common query interchange format, and a wrapper is used to map local query languages (e.g., SQL or XPath) into the common format. Whereas in PeerDB, query reformulation between peers is assisted by agents through a *relation-matching strategy*, a process of matching the metadata of RDBs in different peers.

In **SEWASIE** [6], each peer contains an information node (SINode) that is supported by MOMIS [5], which is a mediator-based system. Within each

SINode, semantic mappings between the local schemas and a *global virtual view* are established in the Global-as-View (GAV) approach. A preliminary discussion on schema mapping representation (also in the GAV approach) using RDF is provided [20], from which the RDFMS language and meta-ontology (for P2P) that we propose in this paper are developed.

3 The Layered Peer Architecture

In a P2P data management system, a peer is supposed to be able to manage its local data source as in a traditional database system. Meanwhile, a peer has to possess the ability of communicating with the other peers by providing and consuming services. For this purpose, we propose for each peer a layered architecture (as shown in Figure 2), by which distributed peers form a pure P2P network.

The peer architecture consists of four layers, in which each upper layer achieves its functionality based on the lower ones. In particular, the *syntax layer* provides a uniform syntax (RDF/XML) for serializing the local ontology and its data, which are converted from the local source residing the peer by a wrapper. The *representation layer* is responsible for representing the local ontology in RDFS and its mappings in RDFMS. The *service layer* implements schema mapping and query processing, which are two main services that a peer can provide to the network. The *application layer* contains a GUI (Graphic User Interface) for interaction with the user which initiates query requests. In summary, the adoption of the layered peer architecture simplifies the resolution of peer-to-peer heterogeneities into level-to-level dependencies, thus facilitating the data interoperation by making the layers more maintainable and reusable. We describe next the key components of the layers.

XML/RDB Wrapper. A wrapper converts the data in the peers into a uniform representation in RDF before they are made accessible to other peers. In our approach, the XML and RDB wrappers are respectively used to transform XML schemas and relational schemas into local ontologies. To facilitate local

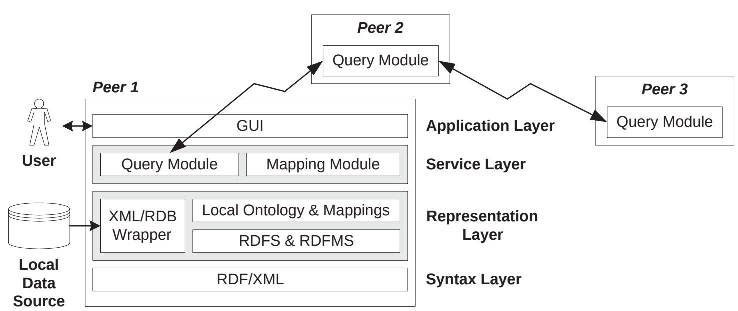

Fig. 2. The layered peer architecture.

query execution, the local data is also converted and serialized (in the syntax layer) and may be managed using a local RDF database (e.g., RSSDB [2]).

Mapping Module. The mapping module is responsible for establishing P2P semantic mappings based on a thesaurus (WordNet[1] in our approach), which is a lexicon database providing semantic clues for concepts and their relationships. More specifically, the mapping module takes as input the two local ontologies (already represented in RDF), finds in the thesaurus a semantically optimal path between each pair of elements in different local ontologies, and derives the direct semantic relationship between elements based on certain rules. This P2P schema matching process and the representation of mappings in RDFMS are discussed in greater detail in Section 4.

Query Module. The query module is responsible for query processing, including query execution, query translation, and answer integration. Through the GUI, the user poses a query, called *local query*, using RQL [12]. The query module then interprets the query and execute it over the *host peer*. Meanwhile, the query is also forwarded to each of the linked peers (called *guest peer*), where the query is rewritten into a *remote query* that is executed locally and propagated further. The query translation utilizes the P2P mapping information established by the mapping module and stored in the representation layer. Finally, the *remote answer* (from the guest peers) and the *local answer* (from the host peer) are returned to the host peer and integrated there as a final answer.

4 RDF-based Metadata Management

4.1 Metadata Representation in RDF

As mentioned previously, to achieve a uniform metadata and data representation in the P2P network, wrappers are used to transform heterogeneous sources into RDF. There exist quite a few methods (e.g., [8, 10, 13]) proposing algorithms for translating data back and forth between different paradigms (e.g., relational, XML, and RDF).

By summarizing existing methods, we choose to convert relations into classes and attributes into properties for translation from an RDB to an RDF-based ontology. For translation from XML to RDF, we convert complex-type elements into classes and simple-type elements (with no subelements but character contents) and attributes into properties. Figure 3 shows the results of schema translation of the three data sources in the example of Figure 1. For the later thesaurus-based schema matching, the naming in the local ontologies is *standardized*, e.g., from `ptitle` to `title`.

In the instance (data) translation that is based on the schema (metadata) translation, important constraints of the local source should be preserved. For example, *foreign key dependencies* between two relations can be translated as a property shared by two classes (corresponding to the two relations) in the target RDF instance. *Nesting constraints* between XML elements can also be preserved

[1] http://www.cogsci.princeton.edu/~wn/

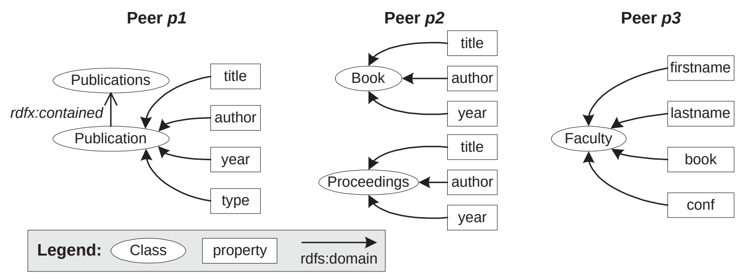

Fig. 3. The RDF-based local ontologies in peers.

by using a new user-defined metaproperty `rdfx:contained` (`rdfx` stands for the namespace), which connects two RDF classes representing the two XML elements [10].

4.2 Thesaurus-based P2P Schema Matching

Schema matching is a long-standing issue in the database community since the 80's. Rahm *et al.* have presented a comprehensive survey of exiting schema matching approaches [19]. Inspired by the work of others [4], we propose in this paper a thesaurus-based schema matching approach for P2P data integration, which utilizes a thesaurus (i.e., WordNet) for discovery of semantic relationships between any two concepts. WordNet is a lexicon database, in which terms are organized into a graph with nodes being a sense (a term may have multiple senses, each being a *synset*, i.e., a set of synonyms) and edges being semantic relations between senses. The WordNet-based schema matching consists of the following three steps (shown in Figure 4).

1. Path Exploration. Among the semantic relations between synsets in WordNet, we choose those of *synonymy*, *hyponymy*/*hypernymy* (i.e., more specific/more general), and *related-to*, when enumerating the paths between two arbitrary but standardized concepts from different local ontologies in peers. In our running example, six paths are found from `Quantity` to `Number`.

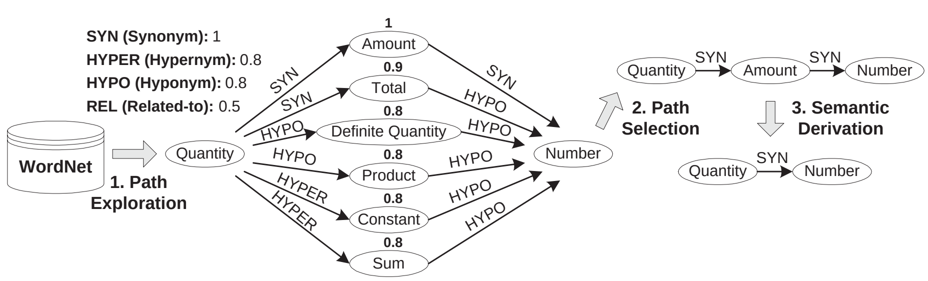

Fig. 4. An example of schema matching using WordNet.

2. Path Selection. When multiple paths are found between two concepts, we choose the *optimal path*, which is supposed to reflect the most likely semantic relation between the two concepts. For this purpose, *semantic similarities* (i.e., the number above each path in the figure) are calculated for all the paths. The calculation is implemented by assigning different semantic relations with different *weights* (e.g., 1 for synonymy and 0.8 for hypernymy, etc.) and then taking the average of all the weights. The path with highest similarity is then chosen as the optimal path. If there is more than one such path, then the user's intervention is needed.

3. Semantic Derivation. The last step is to derive the (direct) semantic relationship, Sem, between the two concepts by reasoning on the semantic relations along the optimal path p between them. More specifically, $Sem(p) = Sem(p_n)$ is computed based on the following recursive algorithm, where $p_n = (r_1, r_2, ..., r_n)$, and $r_i (1 \leq i \leq n)$ are the edges (semantic relations) along p.

$$Sem(p_n) = Sem(p_{n-1}) \wedge Sem(r_n), \qquad if\ n > 1; \qquad (1)$$

$$Sem(p_n) = \quad \approx,\ \supseteq,\ \subseteq,\ or\ \sim, \qquad if\ n = 1. \qquad (2)$$

In the above formulas, the symbols $\approx$, $\supseteq$, $\subseteq$, and $\sim$ respectively stand for semantic relation of synonymy, hypernymy, hyponymy, and related-to. The operation $\wedge$ obeys the rules that are shown in Figure 5, in which a white cell (at the intersection of each pair of gray cells) contains the result of operation on the relations in the two gray cells. Note that the question marks mean that human interference is needed.

The above three steps on each pair of elements in two local ontologies to be mapped enable the determination of the inter-schema mappings. Figure 7 (whose detailed explanation is given in next subsection) gives the resulting mappings among the three ontologies in Figure 3. In this paper we illustrate the core idea of utilizing WordNet for schema matching, while leaving out the detailed path search strategy and its time complexity for future work.

4.3 RDFMS – A Mapping Language

In this subsection, we discuss the issues related to mapping *types* (i.e., semantics), mapping *cardinality* (e.g., one-to-one or one-to-many mappings), and mapping

$\wedge$	$\approx$	$\supseteq$	$\subseteq$	$\sim$
$\approx$	$\approx$	$\supseteq$	$\subseteq$	$\sim$
$\supseteq$	$\supseteq$	$\supseteq$	?	$\sim$
$\subseteq$	$\subseteq$	?	$\subseteq$	$\sim$
$\sim$	$\sim$	$\sim$	$\sim$	$\sim$

Fig. 5. Inference rules on semantic relations.

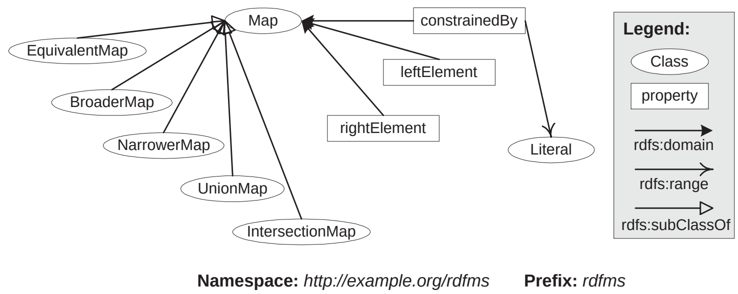

Fig. 6. The meta-ontology of RDFMS.

representation. Every P2P ontology mapping is an instance of the RDFMS meta-ontology.

RDFMS Meta-ontology. To represent the different semantic relationships between concepts in different local ontologies, RDFMS provides such types of mappings as `EquivalentMap` (for synonymy and related-to relations), `BroaderMap` (for hypernymy relations), and `NarrowerMap` (for hyponymy relations). Regarding the case of one-to-many mappings, RDFMS defines `UnionMap` and `IntersectionMap` respectively for two types of logic combinations of the elements on the multiple-element side. Figure 6 shows the meta-ontology represented in the RDF graph model.

In the RDFMS meta-ontology, all types of mappings are defined as classes inheriting from a common class `Map`, which has three general properties that are also inherited by its subclasses. The `leftElement` and `rightElement` properties are used to connect the mapped elements. In order to represent the mapping expression [19] that a P2P mapping may carry, the property `constrainedBy` is defined, whose data type is specified as `Literal`. Following the example in Figure 3, we have the P2P mappings among the three local ontologies in different peers represented in RDFMS, as shown in Figure 7.

Mapping Atoms. In order to use the mapping information for reasoning on mapping and querying, we define a set of mapping atoms as follows.

– $\mathsf{EM}(c_1, c_2)$ iff there exists an instance m of `EquivalentMap`, such that $c_1 = m.$`leftElement` and $c_2 = m.$`rightElement`.
– $\mathsf{BM}(c_1, c_2)$ iff there exists an instance m of `BroaderMap`, such that $c_1 = m.$`leftElement` and $c_2 = m.$`rightElement`.
– $\mathsf{NM}(c_1, c_2)$ iff there exists an instance m of `NarrowerMap`, such that $c_1 = m.$`leftElement` and $c_2 = m.$`rightElement`.
– $\mathsf{UM}(c_1, c_2)$ iff there exists an instance m of `UnionMap`, such that $c_1 = m.$`leftElement` and $c_2 = \{x|x = m.$`rightElement`$\}$, or $c_1 = \{x|x = m.$`rightElement`$\}$ and $c_2 = m.$`leftElement`.
– $\mathsf{IM}(c_1, c_2)$ iff there exists an instance m of `IntersectionMap`, such that $c_1 = m.$`leftElement` and $c_2 = \{x|x = m.$`rightElement`$\}$, or $c_1 = \{x|x = m.$`rightElement`$\}$ and $c_2 = m.$`leftElement`.

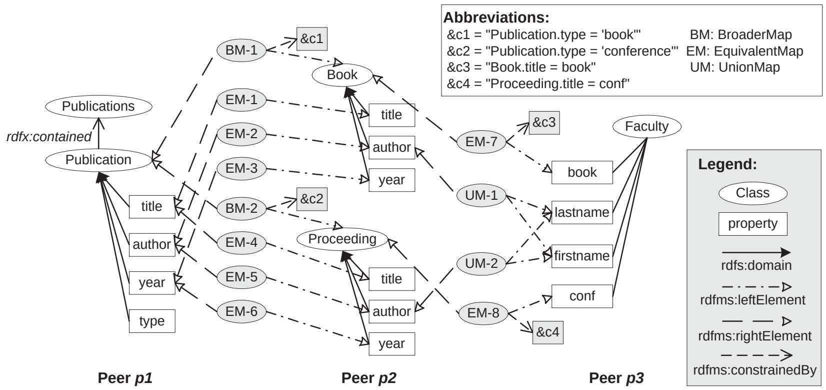

Fig. 7. An example of P2P mappings in RDFMS.

– $\mathsf{CON}(m, e)$ iff given an instance m of `Map` or its subclasses, $e = \{x | x = m.\texttt{constrainedBy}\}$.

Based on these atoms, we can derive new mapping types (e.g., *sibling map* or *disjoint map*) and more complex mappings (e.g., many-to-many mappings by combining two `UnionMaps`). For example, for the new mapping type of *sibling*, SM, we could define such an atom as

$$\mathsf{SM}(c_1, c_2) \equiv \mathsf{NM}(c_1, c_3) \wedge \mathsf{NM}(c_1, c_3).$$

Other Mapping Languages. RDFT [18] is proposed also as a meta-ontology for mapping representation, which is ad-hoc w.r.t the case of business integration. Lehti *et al.* propose an OWL-based model particularly for XML data integration [14]. It uses for mapping representation OWL constructs[1] such as `sameClassAs`, `samePropertyAs`, `subClassOf`, and `subPropertyOf`, thus having a limited power.

5 Queries across the P2P Network

In the syntax layer of our layered peer architecture, we may use an RDF repository (e.g., RSSDB [2]) to store the data converted from local sources. The user then may pose RDF queries over the repository using native languages (e.g., RQL). The basic procedure of query processing (implemented in the query module of a peer) across the P2P network has been described in Section 3. In this section, we focus on the implementation of query translation from peer to peer.

RQL. RQL is a typed functional RDF query language based on a formal model for directed labeled graphs, which permits the interpretation of superimposed resource descriptions by means of one or more RDF schemas [12]. The

[1] http://www.w3.org/2001/sw/WebOnt/

basic form of RQL queries for retrieving RDF data is the SFW expression consisting of `SELECT`, `FROM`, and `WHERE` clauses. For example, the local RQL query Q on Peer p_1 for "listing all publications of a1" may be formulated as follows.

```
SELECT  Y
FROM    {X}title{Y}, {X}author{Z}
WHERE   Z="a1"
```

Most RQL queries are in fact conjunctive queries (called RQL conjunctive queries), which are in the form of $ans(\bar{X}) \; :- \; C_1, ..., C_n$, where C_i's are RQL class or property *patterns* and $\bar{X}$ is a tuple of variables or constants. More details about RQL conjunctive queries can be found in [9].

P2P Query Translation. A first-order relation based model, called SWIM, is also proposed in [9], which we use for the interpretation of RQL conjunctive queries and mapping atoms that are described in Section 4. Because of space limitations, the detailed rules for the interpretation are not discussed here.

Starting from *host peer* p_1, the *local query* (i.e., Q in the example) is propagated across the P2P network. The query module in each peer p_i ($i \neq 1$) is responsible for translating an incoming query into a local query, executing it on the local RDF repository, and forwarding it to all the other linked peers. The query translation can be seen as a function $Q' = f(Q, M)$, where Q, M, and Q' are respectively the incoming query, the mapping information between p_1 and p_i, and the resulting local query. In our approach, the computation of f is simply obtained by composing the underlying SWIM first-order relations of Q with those of M.

Continuing with the running example of this paper, below we show the result of query translation from Q to Q' (in Peer p_2) and to Q'' (in Peer p_3). Note that Q' will be executed over the RDF repository (see Figure 3) in Peer p_2, which contains two RDF schemas. In contrast, Q'' contains two queries, both of which are executed over the RDF data repository in Peer p_3. By integrating all the answers returned from all peers, the final answer to Q is {`"t1", "t3", "t5", "t6"`}.

Q' :
```
SELECT  Y
FROM    {X}title{Y}, {X}author{Z}
WHERE   Z="a1"
```

Q'' :
```
SELECT  Y
FROM    {X}book{Y}, {X}firstname{Z1}, {X}lastname{Z2}
WHERE   Z1="a1" OR Z2="a1"

SELECT  Y
FROM    {X}conf{Y}, {X}firstname{Z1}, {X}lastname{Z2}
WHERE   Z1="a1" OR Z2="a1"
```

6 Conclusions and Future Work

In this paper, an RDF-based layered approach is proposed for data interoperation in a heterogeneous P2P environment. RDF techniques are used throughout the proposed framework, including RDFS for metadata representation, RDFMS (a mapping language and RDF-based meta-ontology) for the representation of P2P mappings, and RQL for query formulation.

At the center of this paper is a thesaurus-based mapping approach that is used to manage metadata in P2P networks. We discuss the process of schema matching, the RDFMS meta-ontology, and the process of reasoning on the mappings. We also describe briefly the query translation process. We note that we could have used OWL to represent the semantic mappings [14], however, the mixture of constructs for the data and the mappings appears to be less suitable for reasons of maintenance, extensibility, and inference.

The approach proposed in this paper is still a first step to a comprehensive solution to P2P data management. A number of research issues remain, including: (1) The application of constraints of the RDFS model and the RDFMS model for query processing. (2) The use of inferencing as applied to mapping and to query processing. (3) The study of the complete syntax and semantics of the mapping language. All these research issues together with the implementation of the framework are our targets in future work.

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