

CS 594 – ADVANCED CRYPTO (SPRING 2026)

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Lecture 1

January 12, 2026

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 - Modern crypto definitions (PPT adversaries)
 - PRGs, PRFs, PRPs
 - MACs
 - OWFs, Random Oracles, Hash Functions
 - PKE (RSA, ElGamal), Key Agreement, Digital Signatures

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- Above is only one part of the field; Crypto is massive!

WHAT IS THIS COURSE?

Tentative list of topics:

- Asymptotic vs. Concrete Security
- Idealized Models
- Secret Sharing
- Zero-knowledge Proofs
- Secure Multi-party Computation
- Indistinguishability Obfuscation
- Memory-hard Functions
- Post-quantum Crypto

SYLLABUS

<https://www.cs.uic.edu/~block/courses/cs594-spring2026/>

MATH REVIEW: NOTATION

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- $\mathbb{N} := \{0, 1, 2, 3, \dots\}$; $\mathbb{Z} := \{0, \pm 1, \pm 2, \pm 3, \dots\}$, $\mathbb{Z}^+ := \{1, 2, 3, \dots\}$,
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- More notation will be added as we progress in the course; please ask questions if things are not clear!

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Definition 1 (Strings)

We let $\{0, 1\}^n$ denote the set of n -bit *binary strings* (i.e., 0/1 vectors of length n). We let $x \in \{0, 1\}^n$ to denote a *binary string* of length n . For any string y , we let $|y|$ denote the *length of y* .

We'll also consider strings over arbitrary alphabets (i.e., sets) S ; e.g., $x \in S^n$ is a string of length n such that $x_i \in S$ for all $i \in \{1, \dots, n\}$.

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- \oplus will be pointwise XOR or addition mod 2
 - $(0, 1) \oplus (1, 1) = (1, 0)$
 - $(0, 1) \oplus (1, 1) = (1, 2) \bmod 2 = (1, 0)$

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- $f(n) = \Theta(g(n))$ if $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$.

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Definition 2 (Negligible Function)

Let $f(n): \mathbb{N} \rightarrow \mathbb{R}_{\geq 0}$. We say that f is *negligible* if for every constant $c > 0$, we have

$$f(n) = o(1/n^c).$$

Ex: 2^{-n} , $\sqrt{2^{-n}} = 2^{-n/2}$
 $q(n) = \text{polynomial}$
 $q(n)/2^n$

MATH REVIEW: EXPONENTS AND LOGARITHMS

Greatest Resource Ever for this: [ChiliMath](#)

$b^{xy} \neq b^x \cdot b^y = b^{x+y}$

$b^x / b^y = b^{x-y}$

$(b^x)^k = b^{x \cdot k}$

$b^0 = 1, b^1 = b$

$b^{-x} = 1/b^x$

$b^{\log_b(k)} = k$

$(ab)^x = a^x b^x$

$\log_b(x \cdot y) = \log_b(x) + \log_b(y)$

$\log_b(x/y) = \log_b(x) - \log_b(y)$

$k \cdot \log_b(x) = \log_b(x^k)$

$\log_b(1) = 0, \log_b(b) = 1$

$-\log_b(x) = \log_b(1/x)$

$\log_b(b^k) = k$

$\log_b(x) = \log_c(x) / \log_c(b)$

↓
change of base

- b is called the *base*
- $x, y, k \in \mathbb{R}$
- for logarithms, the following must hold:
 - $b > 0$ and $b \neq 1$
 - $x, y > 0$

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- $\log := \log_2$ will always denote the binary logarithm (used extensively in this course)

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- Assumption: the base of a logarithm is *always constant*
 - Gives us $\log_b(n) = \Theta(\log_c(n))$ for all constants b, c
 - In other words, doesn't really matter which logarithm base we choose!

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- For $n \in \mathbb{Z}^+$, we let $\mathbb{Z}_n := \{x \bmod n : x \in \mathbb{Z}\}$
 - Abusing notation, $\mathbb{Z}_n \equiv \{0, 1, \dots, n-1\}$, where all operations are performed modulo n

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The *order* of the group is its cardinality $|\mathbb{G}|$. A group is *finite* if the order is finite.

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Definition 4 (Abelian and Cyclic Groups)

A group (\mathbb{G}, \odot) is *Abelian* if \odot is a commutative operator: $a \odot b = b \odot a$ for all $a, b \in \mathbb{G}$. A group (\mathbb{G}, \odot) is *cyclic* if it is Abelian and there exists a *generator* $g \in \mathbb{G}$ such that $\mathbb{G} = \{g^0, g^1, g^2, \dots\}$, where $g^i := \underbrace{g \odot g \odot \dots \odot g}_{i \text{ times}}$, with $g^0 := 1$.

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- In additive groups, we write $i \cdot g = \underbrace{g + g + \dots + g}_{i \text{ times}}$, and
multiplicative groups, we write $g^i = \underbrace{g \cdot g \cdot \dots \cdot g}_{i \text{ times}}$

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without 0

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- For any prime number p , the set $\mathbb{Z}_p^* := \{1, 2, \dots, p-1\}$ is a multiplicative group under multiplication modulo p

MATH REVIEW: FINITE FIELDS

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MATH REVIEW: LINEAR ALGEBRA

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 - For matrices \mathbf{A}, \mathbf{B} , we have $(\mathbf{AB})^\top = \mathbf{B}^\top \cdot \mathbf{A}^\top$ and $(\mathbf{A}^\top)^\top = \mathbf{A}$.

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$$\vec{x} \cdot \vec{y}$$
$$\vec{x} \cdot (\vec{y})^T$$

\mathbb{F} -mult

\mathbb{F} -add

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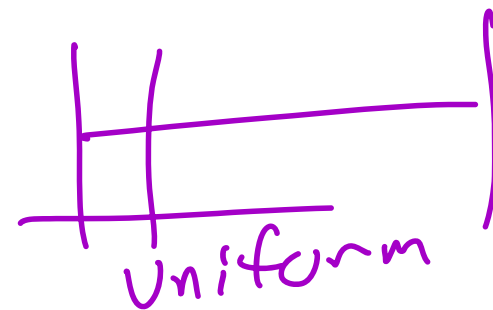
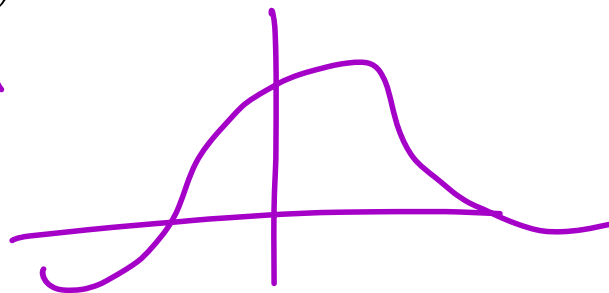
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normal



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 - Note that $\Pr[E] := \sum_{e \in E} \Pr_{s \leftarrow \mathcal{D}} [e = s]$

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↙ conditional probability

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 - **Union Bound:** $\Pr[E_1 \vee E_2] \leq \Pr[E_1] + \Pr[E_2]$

$$\Pr \left[\bigvee_{i=1}^n E_i \right] \leq \sum_{i=1}^n \Pr[E_i]$$

Concrete Security vs. Asymptotic Security